# Iterative Forward Search: Combining Local Search with Maintaining Arc Consistency and a Conflict-based Statistics

### Tomáš Müller<sup>1</sup>,Roman Barták<sup>1</sup>, Hana Rudová<sup>2</sup>

- <sup>1</sup> Faculty of Mathematics and Physics, Charles University {muller|bartak}@ktiml.mff.cuni.cz
- <sup>2</sup> Faculty of Informatics, Masaryk University hanka@fi.muni.cz

### Agenda

- Iterative Forward Search Algorithm
  - Extensions:
    - Conflict-based statistics
    - Maintaining arc consistency
    - Dynamic backtracking
- Experiments
  - Purdue University Timetabling Problem
  - Random binary CSP
- Conclusion

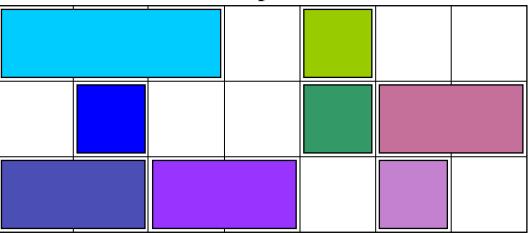


#### **Iterative Forward Search**

- Basic Approaches
  - Local search
  - Backtracking based search
- Iterative Forward Search Algorithm
  - Forward based search
  - Works in iterations
  - Extending a (partial) feasible solution
  - Interactivity



A (partial) feasible solution

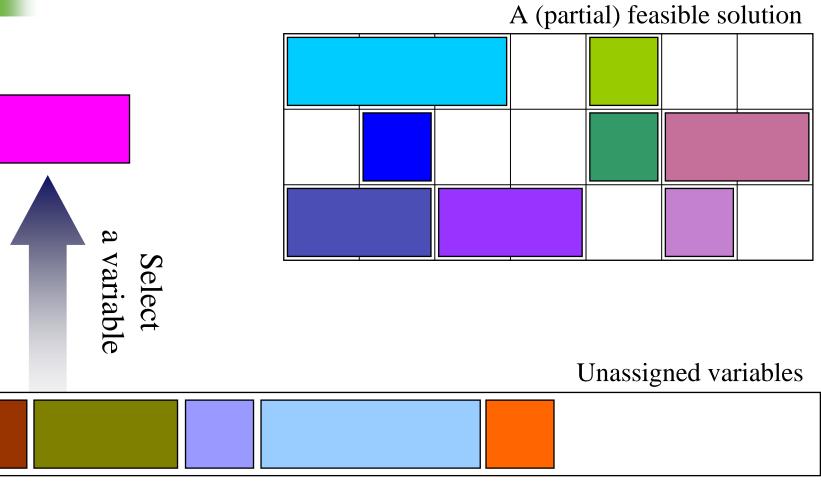


Unassigned variables

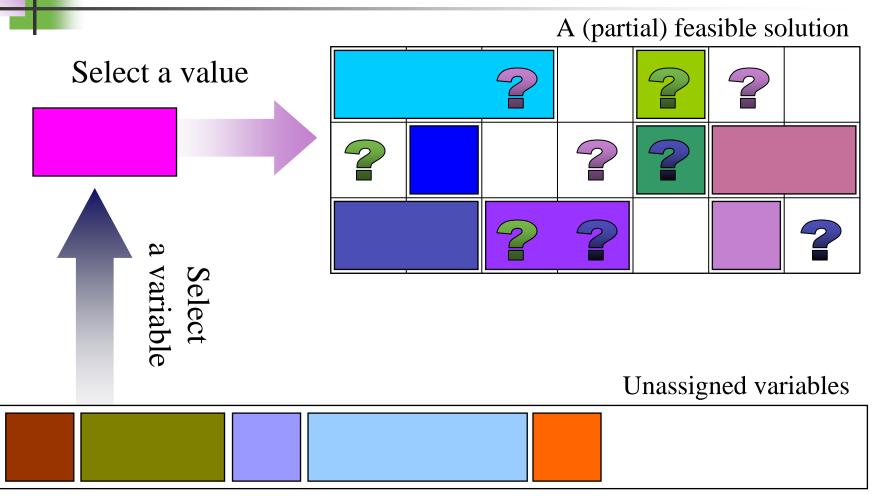


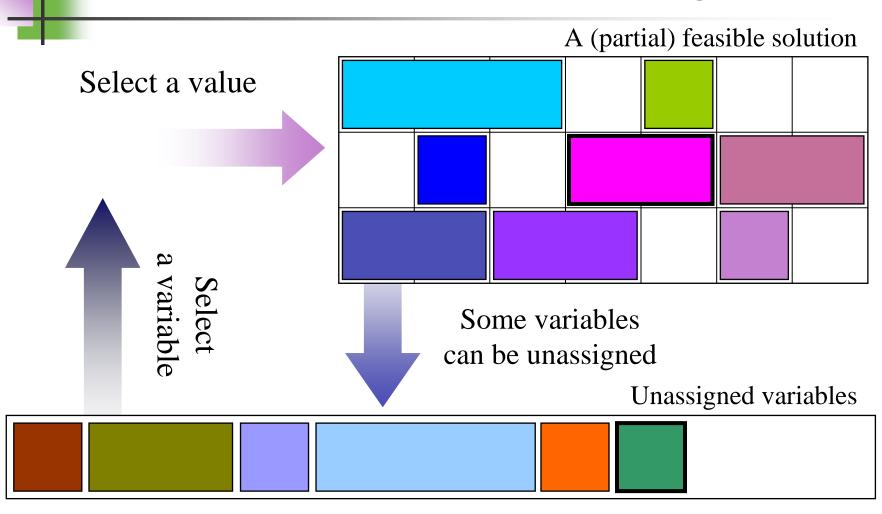
### Iterative For

Iterative Forward Search Algorithm











### Guided by

- Variable selection
  - First-fail principle
- Value selection
  - Best-fit value
- Solution comparator
  - Less unassigned variables, less violated soft constraints, ...
- Termination condition
  - Solution is complete and good enough
  - Timeout or user interaction

### Conflict-based statistics

- Idea
  - Memorize conflicts and discourage their potential repetition
- If B=c is unassigned because of the A=a
  - A counter Stat[ $A=a,B\neq c$ ] is incremented

$$A = a \Rightarrow \begin{cases} 3 \times B \neq a \\ 4 \times B \neq c \\ 2 \times C \neq a \\ 120 \times D \neq a \end{cases}$$



### Conflict-based statistics

### To be used e.g. in value selection

- If a is being selected for variable A
- And there is B=b in a conflict with A=a



Value a is weighted by  $Stat[A=a,B\neq b]+1$ 

Conflicts are weighted by their occurrences in the past

### 1

### Maintaining Arc Consistency

#### Based on explanations

- $V_i \neq v_i \Leftarrow (V_1 = v_1 \& V_2 = v_2 \dots \& V_j = v_j)$
- When a value is removed from a domain
  - An explanation is attached to the deleted value
- When a variable is unassigned (e.g.,  $V_x = v_x$ )
  - All deleted values which contain  $V_x = v_x$  in their explanations have to be recomputed
- Computation
  - FC: Explanation corresponds to the violated constraint
  - MAC: Union of explanations

### 1

### Dynamic Backtracking with MAC

- A special case of IFS with MAC
  - An unassigned variable is always selected
  - If there is a variable with an empty domain
    - A union of assignments of all values' explanations is computed
    - Fail if the computed union is empty
    - The last assignment from the union is unassigned
    - Explanation: all the other assignments in the computed union
  - If a value  $v_x$  is assigned to  $V_x$ 
    - An explanation  $V_x \neq v'_x \leftarrow (V_x = v_x)$  is attached to all values from the variables domain different from  $v_x$



- Timetabling Problem at Purdue University
  - Central timetable for large lecture classes
    - 826 classes (forming 1782 meetings)

Fall 2004

- some of them (25%) with multiple sections
- 50 lecture rooms (with various equipment, up to 474 seats)
- 89,633 course demands from 29,808 students
- Utilization over 78% (~ 94% for the four largest rooms)
- Timetables for individual departments
  - Done manually for the moment
    - An area for our future work

### Experiments: Purdue University Timetabling

#### For each class

- Student requirements
- Time requirements & preferences
  - Meeting patterns (e.g., 3 x 50 min, 2 x 75 min)
- Room requirements & preferences
  - Capacity
  - Required equipment
  - Room / building preference
- Instructor
- Additional (group) constraints
  - Between several classes (e.g. back-to-back, precedence)
- Other ...

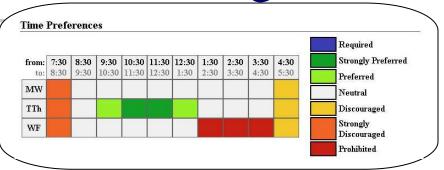
Each student states which courses he or she wants to attend (soft constraint)

### Experiments:

Purdue University Timetabling



- Student requirements
- Time requirements & preferences/
  - Meeting patterns (e.g., 3 x 50 min, 2 x 75 min)
- Room requirements & preferences
  - Capacity
  - Required equipment
  - Room / building preference
- Instructor
- Additional (group) constraints
  - Between several classes (e.g. back-to-back, precedence)
- Other ...



## Experiments: Purdue University Timetabling

#### For each class

- Student requirements
- Time requirements & preferences

Meeting patterns (e.g., 3 x 50 min, 2 x 75 min)

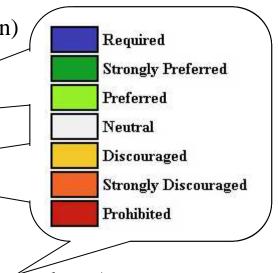
Room requirements & preferences

Capacity

Required equipment

Room / building preference

- Instructor
- Additional (group) constraints
  - Between several classes (e.g. back-to-back, précedence)
- Other ...



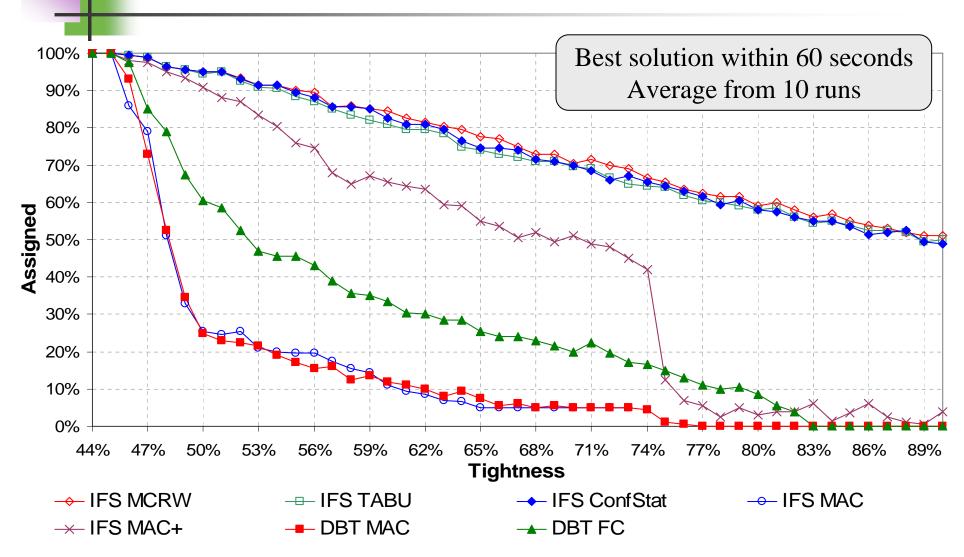
### Experiments: Purdue University Timetabling

Test Case	IFS ConfStat	IFS TABU	IFS MCRW
Assigned classes [%]	$100.0 \pm 0.00$	$97.67 \pm 0.15$	$98.29 \pm 0.16$
Time [min]	$24.11 \pm 4.42$	$24.17 \pm 3.62$	$24.52 \pm 3.83$
Student conflicts [%]	$1.97 \pm 0.06$	$\boldsymbol{1.97} \pm 0.07$	$2.05 \pm 0.19$
Preferred time [%]	$85.64 \pm 1.57$	$89.86 \pm 0.69$	$89.63 \pm 1.06$
Preferred room [%]	$50.39 \pm 5.34$	$66.48 \pm 3.42$	$64.84 \pm 3.86$

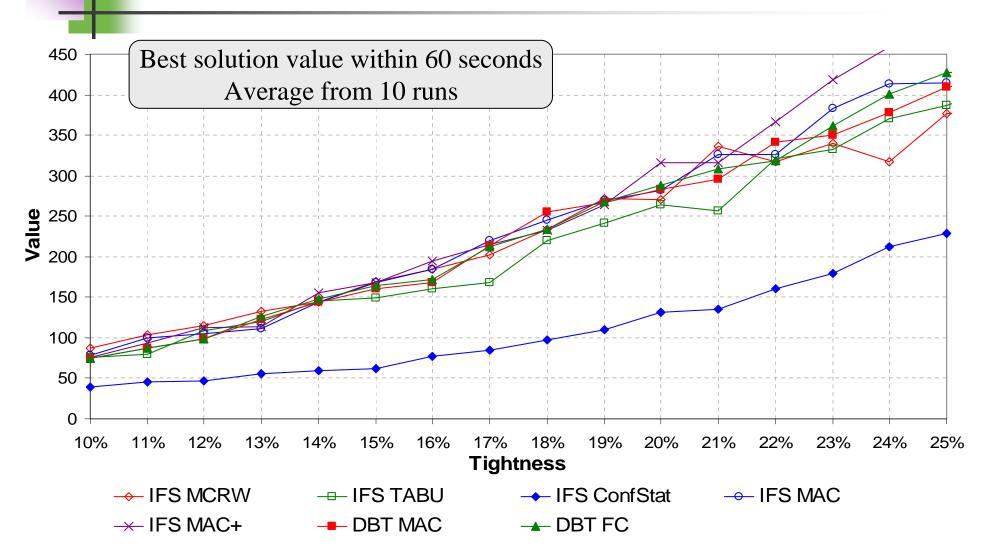
- DBT MAC was able to assign approx. 93% of variables
- IFS MAC was able to assign approx. 94% of variables

Best solution within 30 minutes, 10 runs 1 GHz Pentium III, Java 1.4.2











### Conclusion And Future Work

- IFS algorithm with conflict-based statistics
  - Good results on Purdue University Timetabling Problem
- Future work
  - More results
    - Timetables for individual departments
    - On other (not only timetabling) problems
  - Solver improvements
  - Additional requirements from Purdue University
  - Application of conflict-based statistics in other search techniques